

Compiler I

(dt. Übersetzer I)

Prof. Dr. Uwe Kastens

Winter 2001/2002

Objectives

The participants are taught to

- understand **fundamental techniques** of language implementation,
- use **generating tools and standard solutions**,
- understand compiler construction as a systematic combination of **algorithms, theories** and **software engineering** methods for the solution of a **precisely specified task**,
- apply compiler techniques for languages **other than programming languages**.

Forms of teaching:

Lectures

Tutorials

Homeworks

Exercises

Running project

Lectures in English

Some agreements about giving lectures in English:

- I'll speak English unless someone asks me to explain something in German.
- Stop me or slow me down whenever you get lost.
- I don't speak as well as a native speaker; but I'll do my best ...
- You may ask questions and give answers in English or in German.
- I'll prepare the slides in English. A German version is available.
- You'll have to learn to speak about the material in at least one of the two languages.
- You may vote which language to be used in the tutorials.
- You may chose German or English for the oral exam.

Syllabus

Week	Chapter	Topic
1	Introduction	Compiler tasks
2		Compiler structure
3	Lexical analysis	Scanning, token representation
4	Syntactic analysis	Recursive decent parsing
5		LR Parsing
6		Parser generators
7		Grammar design
8	Semantic analysis	Attribute grammars
9		Attribute grammar specifications
10		Name analysis
11		Type analysis
12	Transformation	Intermediate language, target trees
13		Target texts
14	Synthesis	Overview
15	Summary	

Prerequisites

from Lecture	Topic	here needed for
Foundations of Programming Languages:		
	4 levels of language properties	Compiler tasks, compiler structure
	Context-free grammars	Syntactic analysis
	Scope rules	Name analysis
	Data types	Type analysis
	Lifetime, runtime stack	Storage model, code generation
Modeling:		
	Finite automata	Lexical analysis
	Context-free grammars	Syntactic analysis

References

Material for this course **Compiler I**: <http://www.uni-paderborn.de/cs/ag-kastens/compi>
 in German **Übersetzer I** (1999/2000): <http://www.uni-paderborn.de/cs/ag-kastens/uebi>
 in English **Compiler II**: <http://www.uni-paderborn.de/cs/ag-kastens/uebii>

Modellierung: <http://www.uni-paderborn.de/cs/ag-kastens/model>
Grundlagen der Programmiersprachen: <http://www.uni-paderborn.de/cs/ag-kastens/gdp>

U. Kastens: **Übersetzerbau**, Handbuch der Informatik 3.3, Oldenbourg, 1990
 (not available on the market anymore, available in the library of the University)

W. M. Waite, L. R. Carter: **An Introduction to Compiler Construction**, Harper Collins, New York, 1993

W. M. Waite, G. Goos: **Compiler Construction**, Springer-Verlag, 1983

R. Wilhelm, D. Maurer: **Übersetzerbau - Theorie, Konstruktion, Generierung**, Springer-Verlag, 1992

A. Aho, R. Sethi, J. D. Ullman: **Compilers - Principles, Techniques and Tools**, Addison-Wesley, 1986

A. W. Appel: **Modern Compiler Implementation in C**, Cambridge University Press, 1997
 (available for Java and for ML, too)

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Course material in the Web

Lecture
Compiler I WS 2001/2002
 Prof. Dr. Uwe Kastens
[other lectures](#)

Slides	Organization	Supplements
<ul style="list-style-type: none"> forward / backward Contents Printing 	<ul style="list-style-type: none"> general actual information <p>18.09.2001 First lecture: Monday, Oct 15 18.09.2001 Room change: D1338</p>	<ul style="list-style-type: none"> Objectives Site map Literature Contents Kastens: Übersetzerbau Internet Material in German

Exercises

- [forward / backward](#)
- [Overview](#)
- [Printing](#)

Usage

We recommend to use the full screen size for the browser window. You may even hide the directory buttons to minimize the need for scrolling.

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Commented slide in the course material

Lecture Compiler I WS 2001/2002 – Slide no. 25

Compilation and interpretation of Java programs

Source modules

Java Compiler

Class files in Java Bytecode (intermediate language)

load needed class files dynamically - local or via Internet

Class loader

Bytecode processor in software

Just-In-Time Compiler (JIT)

Machine code

Interpreter Java Virtual Machine JVM

Input

Output

Objectives:
 Special situation for Java

In the lecture:
 Explain the role of the abstract machine JVM:

- Interpretation of bytecode.
- Compile and optimize while executing the program.
- Load class files while executing the program.

Questions:

- explain why the JVM can not rely on the type checks made by the compiler.

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What does a compiler compile?

A **compiler** transforms correct sentences of its **source language** into sentences of its **target language** such that their **meaning is unchanged**.

Examples:

Source language:	Target language:
Programming language	Machine language
C++	Sparc code
Programming language	Abstract machine
Java	Java Bytecode
Programming language	Programming language (source-to-source)
C++	C
Application language	Application language
LaTeX	HTML
Data base language (SQL)	Data base system calls

What is compiled here?

```
class Average
{ private:
  int sum, count;
public:
  Average (void)
  { sum = 0; count = 0; }
  void Enter (int val)
  { sum = sum + val; count++; }
  float GetAverage (void)
  { return sum / count; }
};
```

```
-----
_Enter__7Averagei:
    pushl %ebp
    movl %esp,%ebp
    movl 8(%ebp),%edx
    movl 12(%ebp),%eax
    addl %eax,(%edx)
    incl 4(%edx)

L6:
    movl %ebp,%esp
    popl %ebp
    ret
```

```
class Average
{ private:
  int sum, count;
public:
  Average ()
  { sum = 0; count = 0; }
  void Enter (int val)
  { sum = sum + val; count++; }
  float GetAverage ()
  { return sum / count; }
};

1: Enter: (int) --> void
Access: []
Attribute 'Code' (Length 49)
Code: 21 Bytes Stackdepth: 3 Locals: 2
0:   aload_0
1:   aload_0
2:   getfield cp4
5:   iload_1
6:   iadd
7:   putfield cp4
10:  aload_0
11:  dup
12:  getfield cp3
15:  iconst_1
16:  iadd
```

What is compiled here?

```
program Average;
  var sum, count: integer;
  aver: integer;
  procedure Enter (val: integer);
  begin sum := sum + val;
        count := count + 1;
  end;
begin
  sum := 0; count := 0;
  Enter (5); Enter (7);
  aver := sum div count;
end.

-----
void ENTER_5 (char *slnk , int VAL_4)
{
  /* data definitions: */
  /* executable code: */
  {
    SUM_1 = (SUM_1)+(VAL_4);
    COUNT_2 = (COUNT_2)+(1);
  }
} /* ENTER_5 */
```

```
\documentstyle[12pt]{article}
\begin{document}
\section{Introduction}
This is a very short document.
It just shows
\begin{itemize}
\item an item, and
\item another item.
\end{itemize}
\end{document}

-----
%%Page: 1 1
1 0 bop 164 315 a Fc(1)81
b(In)n(tro)r(duction)
164 425 y Fb(This)16
b(is)g(a)h(v)o(ery)e(short)
i(do)q(cumen)o(t.)j(It)c(just)g
(sho)o(ws)237 527 y Fa(\017)24 b
Fb(an)17 b(item,)
c(and)237 628 y Fa(\017)24 b
Fb(another)17 b(item.)
961 2607 y(1)p
eop
```

Languages for specification and modeling

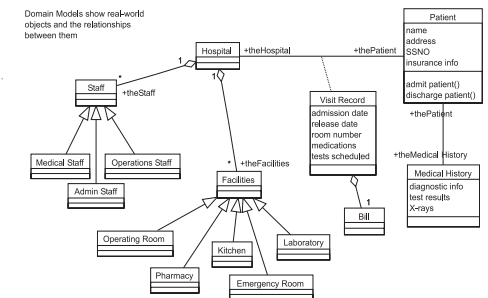
SDL (CCITT)
Specification and Description Language:

```
block Dialogue;
  signal
    Money, Release, Change, Accept, Avail, Unavail, Price,
    Showtxt, Choice, Done, Flushed, Close, Filled;
  process Coins referenced;
  process Control referenced;
  process Viewpoint referenced;
  signalroute Plop
    from env to Coins
      with Coin_10, Coin_50, Coin_100, Coin_x;
  signalroute Pong
    from Coins to env
      with Coin_10, Coin_50, Coin_100, Coin_x;
  signalroute Cash
    from Coins to Control
      with Money, Avail, Unavail, Flushed, Filled;
  from Control to Coins
    with Accept, Release, Change, Close;

...

connect Pay and Plop;
connect Flush and Pong;
endblock Dialogue;
```

UML
Unified Modeling Language:



Domain Specific Languages (DSL)

CI-12

A language designed for a **specific application domain**.

Application Generator: Implementation of a DSL by a **program generator**

Examples:

- **Simulation of mechatronic feedback systems**
- **Robot control**
- **Collecting data from instruments**
- **Testing car instruments**
- **Report generator for bibliographies:**

```
string name =   InString "Which author?";
int since =     InInt "Since which year?";
int cnt = 0;

"\nPapers of ", name, " since ", since, ":\n";
[ SELECT name <= Author && since <= Year;
  cnt = cnt + 1;
  Year, "\t", Title, "\n";
]
"\n", name, " published ", cnt, "papers.\n";
```

U. Kastens: Construction of
Application Generators
Using Eli,
Workshop on Compiler
Techniques for Application
Domain Languages ...,
Linköping, April 1996

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Programming languages as source or target languages

CI-13

Programming languages as source languages:

- **Program analysis**
call graphs, control-flow graph, data dependencies, e. g. for the year 2000 problem
- **Recognition of structures and patterns**
e. g. for Reengineering

Program languages as target languages:

- **Specifications (SDL, OMT, UML)**
- **graphic modeling of structures**
- **DSL, Application generator**

=> Compiler task: Source-to-source compilation

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Semester project as running example

CI-14

A Structure Generator

We are going to develop a tool that implements **record structures**. In particular, the structure generator takes a set of **record descriptions**. Each specifies a **set of named and typed fields**. For each record a **Java class** declaration is to be generated. It contains a constructor method and access methods for the specified record fields.

The tool will be used in an environment where field description are created by other tools, which for example analyze texts for the occurrence of certain phrases. Hence, the descriptions of fields may occur in arbitrary order, and the same field may be described more than once. The structure generator **accumulates the field descriptions** such that for each record a single class declaration is generated which has all fields of that record.

Design a **domain specific language**.

Implement an **application generator** for it.

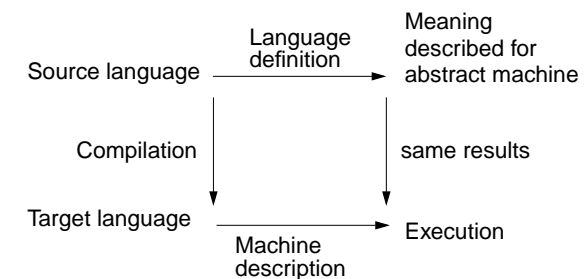
Apply all **techniques of the course** that are useful for the task.

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Meaning preserving transformation

CI-15

A **compiler** transforms correct sentences of its **source language** into sentences of its **target language** such that their **meaning is unchanged**.



A **meaning** is defined only for **correct** programs. Compiler task: Error handling

The compiler analyses **static** properties of the program at **compile time**, e. g. definitions of Variables, types of expressions. Decides: Is the program **compilable**?

Dynamic properties of the program are checked at **runtime**, e. g. indexing of arrays. Decides: Is the program **executable**?

But in Java: Compilation of bytecode at runtime, just in time compilation (JIT)

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Example: Tokens and structure

Character sequence

```
int count = 0; double sum = 0.0; while (count < maxVect) { sum = sum + vect[count]; count++; }
```

Tokens

```
int count = 0; double sum = 0.0; while (count < maxVect) { sum = sum + vect[count]; count++; }
```

Expressions

Declarations

Statements

Structure

Example: Names, types, generated code

Tokens

```
int count = 0; double sum = 0.0; while (count < maxVect) { sum = sum + vect[count]; count++; }
```



k1: (count, local variable, int)
k2: (sum, local variable, double)

k3: (maxVect, member variable, int) ...
k4: (vect, member variable, double array)

Names and types

generated Bytecode

```
0 iconst_0          12 faload
1 istore_1          13 f2d
2 dconst_0          14 dadd
3 dstore_2          15 dstore_2
4 goto 19           16 iinc 1 1
7 dload_2           19 iload_1
8 getstatic #5 <vect[]> 20 getstatic #4 <maxVect>
11 iload_1          23 if_icmplt 7
```

Language definition - Compiler task

• Notation of tokens

keywords, identifiers, literals
formal definition: regular expressions

lexical analysis

• Syntactic structure

formal definition: context-free grammar

syntactic analysis

• Static semantics

binding names to program objects, typing rules
usually defined by informal texts

semantic analysis, transformation

• Dynamic semantics

semantics, effect of the execution of constructs
usually defined by informal texts
in terms of an abstract machine

transformation, code generation

• Definition of the target language (machine)

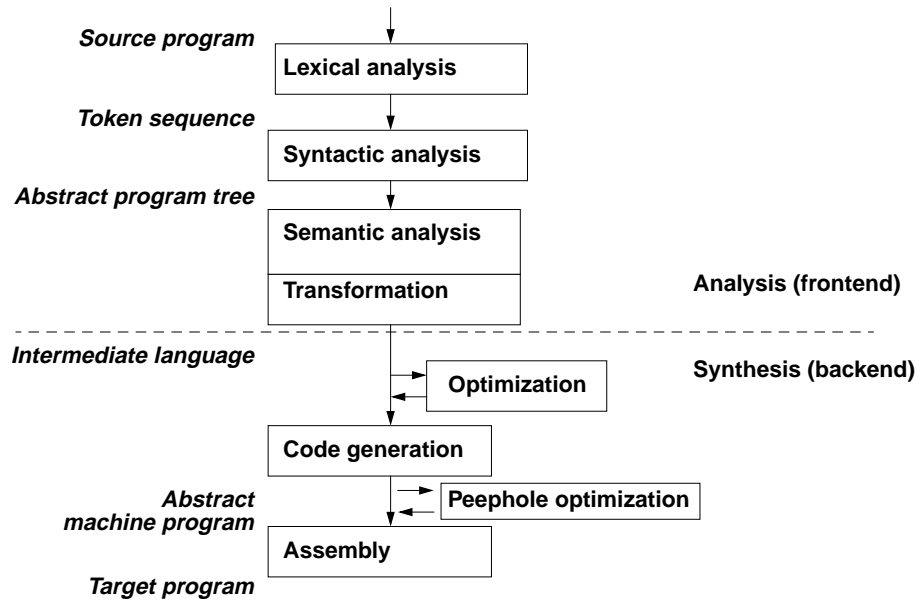
transformation, code generation assembly

Compiler tasks

Structuring	Lexical analysis	Scanning Conversion
	Syntactic analysis	Parsing Tree construction
Translation	Semantic analysis	Name analysis Type analysis
	Transformation	Data mapping Action mapping
Encoding	Code generation	Execution-order Register allocation Instruction selection
	Assembly	Instruction encoding Internal Addressing External Addressing

Compiler structure and interfaces

CI-20



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Software qualities of the compiler

CI-21

- **Correctness** Translate correct programs correctly.
Reject wrong programs and give error messages
- **Efficiency** Storage and time used by the compiler
- **Code efficiency** Storage and time used by the generated code
Compiler task: Optimization
- **User support** Compiler task: Error handling
(recognition, message, recovery)
- **Robustness** Give a reasonable reaction on every input

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Strategies for compiler construction

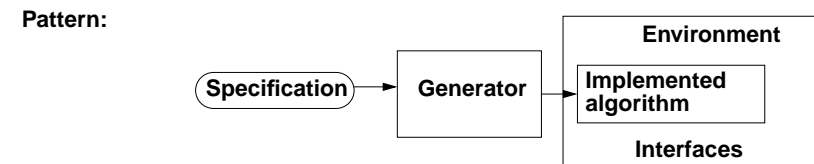
CI-22

- Obey exactly to the language definition
- Use generating tools
- Use standard components
- Apply standard methods
- Validate the compiler against a test suite
- Verify components of the compiler

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Generators

CI-23



Typical compiler tasks solved by generators:

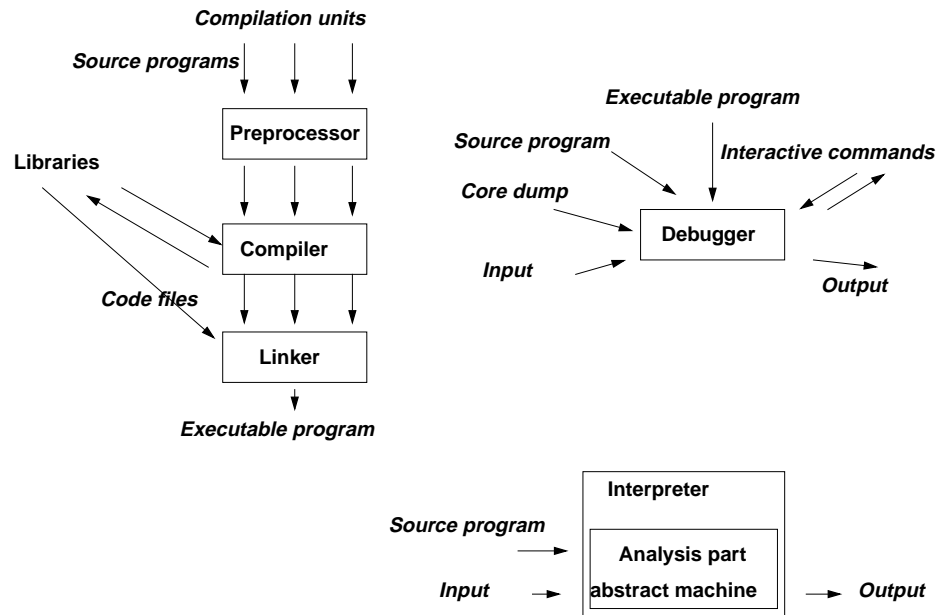
Regular expressions	Scanner generator	Finite automaton
Context-free grammar	Parser generator	Stack automaton
Attribute grammar	Attribute evaluator generator	Tree walking algorithm
Code patterns	Code selection generator	Pattern matching

integrated system Eli:



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Environment of compilers



Compilation and interpretation of Java programs

