

# **Compiler I**

## **(dt. Übersetzer I)**

**Prof. Dr. Uwe Kastens**

**Winter 2001/2002**

# Objectives

The participants are taught to

- understand **fundamental techniques** of language implementation,
- use **generating tools and standard solutions**,
- understand compiler construction as a systematic combination of **algorithms, theories** and **software engineering** methods for the solution of a **precisely specified task**,
- apply compiler techniques for languages **other than programming languages**.

Forms of teaching:

**Lectures**

**Tutorials**

**Homeworks**

**Exercises**

**Running project**

# Lectures in English

Some agreements about giving lectures in English:

- I'll speak English unless someone asks me to explain something in German.
- Stop me or slow me down whenever you get lost.
- I don't speak as well as a native speaker; but I'll do my best ...
- You may ask questions and give answers in English or in German.
- I'll prepare the slides in English. A German version is available.
- You'll have to learn to speak about the material in at least one of the two languages.
- You may vote which language to be used in the tutorials.
- You may chose German or English for the oral exam.

# Syllabus

Week	Chapter	Topic
1	<b>Introduction</b>	Compiler tasks
2		Compiler structure
3	<b>Lexical analysis</b>	Scanning, token representation
4	<b>Syntactic analysis</b>	Recursive decent parsing
5		LR Parsing
6		Parser generators
7		Grammar design
8	<b>Semantic analysis</b>	Attribute grammars
9		Attribute grammar specifications
10		Name analysis
11		Type analysis
12	<b>Transformation</b>	Intermediate language, target trees
13		Target texts
14	<b>Synthesis</b>	Overview
15	<b>Summary</b>	

# Prerequisites

**from Lecture    Topic**

**here needed for**

Foundations of Programming Languages:

4 levels of language properties

Compiler tasks, compiler structure

Context-free grammars

Syntactic analysis

Scope rules

Name analysis

Data types

Type analysis

Lifetime, runtime stack

Storage model, code generation

Modeling:

Finite automata

Lexical analysis

Context-free grammars

Syntactic analysis

# References

Material for this course **Compiler I:** <http://www.uni-paderborn.de/cs/ag-kastens/compi>  
in German **Übersetzer I** (1999/2000): <http://www.uni-paderborn.de/cs/ag-kastens/uebi>  
in English **Compiler II:** <http://www.uni-paderborn.de/cs/ag-kastens/uebii>

**Modellierung:** <http://www.uni-paderborn.de/cs/ag-kastens/model>  
**Grundlagen der Programmiersprachen:** <http://www.uni-paderborn.de/cs/ag-kastens/gdp>

U. Kastens: **Übersetzerbau**, Handbuch der Informatik 3.3, Oldenbourg, 1990  
(not available on the market anymore, available in the library of the University)

W. M. Waite, L. R. Carter: **An Introduction to Compiler Construction**,  
Harper Collins, New York, 1993

W. M. Waite, G. Goos: **Compiler Construction**, Springer-Verlag, 1983

R. Wilhelm, D. Maurer: **Übersetzerbau - Theorie, Konstruktion, Generierung**,  
Springer-Verlag, 1992

A. Aho, R. Sethi, J. D. Ullman: **Compilers - Principles, Techniques and Tools**,  
Addison-Wesley, 1986

A. W. Appel: **Modern Compiler Implementation in C**, Cambridge University Press, 1997  
(available for Java and for ML, too)

# Course material in the Web

Netscape: Lecture Compiler I WS 2001/2002

File Edit View Go Communicator Help

Back Forward Reload Home Search Netscape Print Security Shop Stop


Bookmarks Location: <http://www.uni-paderborn.de/cs/ag-kastens/compi/>

**Lecture**

**Compiler I WS 2001/2002**

**Prof. Dr. Uwe Kastens**

[other lectures](#)

  
Universität Paderborn  
Praktische Informatik

Slides	Organization	Supplements
<ul style="list-style-type: none"> <li>• <a href="#">forward / backward</a></li> <li>• <a href="#">Contents</a></li> <li>• <a href="#">Printing</a></li> </ul>	<ul style="list-style-type: none"> <li>• <a href="#">general</a></li> <li>• <a href="#">actual information</a></li> </ul> <p>18.09.2001 <a href="#">First lecture: Monday, Oct 15</a></p> <p>18.09.2001 <a href="#">Room change: D1.338</a></p>	<ul style="list-style-type: none"> <li>• <a href="#">Objectives</a></li> <li>• <a href="#">Site map</a></li> <li>• <a href="#">Literature</a></li> <li>• <a href="#">Contents Kastens: Übersetzerbau</a></li> <li>• <a href="#">Internet</a></li> <li>• <a href="#">Material in German</a></li> </ul>

**Exercises**

- [forward / backward](#)
- [Overview](#)
- [Printing](#)

**Usage**

We recommend to use the full screen size for the browser window. You may even hide the directory buttons to minimize the need for scrolling.

This material is maintained by **CAMELOT**.

# Commented slide in the course material

Netscape: Lecture Compiler I WS 2001/2002 / Slide 25

File Edit View Go Communicator Help

Back Forward Reload Home Search Netscape Print Security Shop Stop

Bookmarks Location: <http://www.uni-paderborn.de/cs/ag-kastens/compi/folien/Folie25.html>

## Lecture Compiler I WS 2001/2002 – Slide no. 25

CI-25

### Compilation and interpretation of Java programs

```

graph TD
    SM[Source modules] --> JC[Java Compiler]
    JC --> CFB[Class files in Java Bytecode intermediate language]
    CFB --> CL[Class loader]
    CL --> BPP[Bytecode processor in software]
    BPP --> JIT[Just-In-Time Compiler JIT]
    JIT --> MC([Machine code])
    MC --> Output
    Input --> CL
    CL --> CLD[load needed class files dynamically - local or via Internet]
    CLD --> CL
  
```

The diagram illustrates the flow of Java programs from source modules to machine code. Source modules are compiled by the Java Compiler into class files in Java Bytecode (intermediate language). These class files are then loaded by the Class loader, which can load needed class files dynamically, either locally or via the Internet. The loaded class files are then processed by the Bytecode processor in software, which then feeds into the Just-In-Time Compiler (JIT). The JIT compiler produces machine code, which is the final output. The entire process from class files to machine code is part of the Interpreter Java Virtual Machine (JVM).

Objectives:  
Special situation for Java

In the lecture:  
Explain the role of the abstract machine JVM:

- Interpretation of bytecode.
- Compile and optimize while executing the program.
- Load class files while executing the program.

Questions:

- explain why the JVM can not rely on the type checks made by the compiler.

© 2001 bei Prof. Dr. Uwe Kastens



# What does a compiler compile?

A **compiler** transforms correct sentences of its **source language** into sentences of its **target language** such that their **meaning is unchanged**.

## Examples:

### Source language:

#### Programming language

C++

#### Programming language

Java

#### Programming language

C++

#### Application language

LaTeX

Data base language (SQL)

### Target language:

#### Machine language

Sparc code

#### Abstract machine

Java Bytecode

#### Programming language (source-to-source)

C

#### Application language

HTML

Data base system calls

# What is compiled here?

```
class Average
{ private:
    int sum, count;
public:
    Average (void)
        { sum = 0; count = 0; }
    void Enter (int val)
        { sum = sum + val; count++; }
    float GetAverage (void)
        { return sum / count; }
};
```

-----  
\_Enter\_\_7Averagei:

```
    pushl %ebp
    movl %esp,%ebp
    movl 8(%ebp),%edx
    movl 12(%ebp),%eax
    addl %eax, (%edx)
    incl 4(%edx)

L6:
    movl %ebp,%esp
    popl %ebp
    ret
```

```
class Average
{ private
    int sum, count;
public
    Average ()
        { sum = 0; count = 0; }
    void Enter (int val)
        { sum = sum + val; count++; }
    float GetAverage ()
        { return sum / count; }
};
```

-----  
1: Enter: (int) --> void

Access: []

Attribute ,Code` (Length 49)

Code: 21 Bytes Stackdepth: 3 Locals: 2

```
0:      aload_0
1:      aload_0
2:      getfield cp4
5:      iload_1
6:      iadd
7:      putfield cp4
10:     aload_0
11:     dup
12:     getfield cp3
15:     iconst_1
16:     iadd
```

# What is compiled here?

```

program Average;
  var sum, count: integer;
      aver: integer;
  procedure Enter (val: integer);
    begin sum := sum + val;
          count := count + 1;
    end;
begin
  sum := 0; count := 0;
  Enter (5); Enter (7);
  aver := sum div count;
end.

```

```

-----
void ENTER_5 (char *slnk , int VAL_4)
{
  /* data definitions: */
  /* executable code: */
  {
    SUM_1 = (SUM_1)+(VAL_4);
    COUNT_2 = (COUNT_2)+(1);
    ;
  }
} /* ENTER_5 */

```

```

\documentstyle[12pt]{article}
\begin{document}
\section{Introduction}
This is a very short document.
It just shows
\begin{itemize}
\item an item, and
\item another item.
\end{itemize}
\end{document}

```

-----

```

%%Page: 1 1
1 0 bop 164 315 a Fc(1)81
b(In)n(tro)r(duction)
164 425 y Fb(This)16
b(is)g(a)h(v)o(ery)e(short)
i(do)q(cumen)o(t.)j(It)c(just)g
(sho)o(ws)237 527 y Fa(\017)24 b
Fb(an)17 b(item,)
c(and)237 628 y Fa(\017)24 b
Fb(another)17 b(item.)
961 2607 y(1)p
eop

```

# Languages for specification and modeling

## SDL (CCITT)

Specification and Description Language:

```

block Dialogue;
  signal
    Money, Release, Change, Accept, Avail, Unavail, Price,
    Showtxt, Choice, Done, Flushed, Close, Filled;
  process Coins referenced;
  process Control referenced;
  process Viewpoint referenced;
  signalroute Plop
    from env to Coins
      with Coin_10, Coin_50, Coin_100, Coin_x;
  signalroute Pong
    from Coins to env
      with Coin_10, Coin_50, Coin_100, Coin_x;
  signalroute Cash
    from Coins to Control
      with Money, Avail, Unavail, Flushed, Filled;
    from Control to Coins
      with Accept, Release, Change, Close;

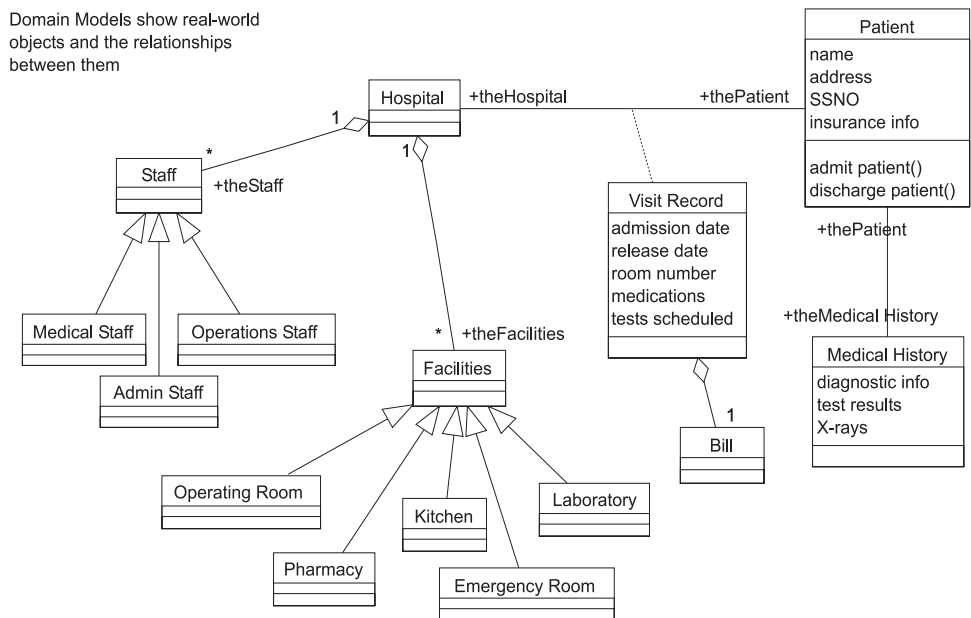
  ...

  connect Pay and Plop;
  connect Flush and Pong;
endblock Dialogue;
  
```

## UML

Unified Modeling Language:

Domain Models show real-world objects and the relationships between them



# Domain Specific Languages (DSL)

A language designed for a **specific application domain**.

**Application Generator:** Implementation of a DSL by a **program generator**

## Examples:

- **Simulation of mechatronic feedback systems**
- **Robot control**
- **Collecting data from instruments**
- **Testing car instruments**
- **Report generator for bibliographies:**

```
string name =   InString "Which author?";
int since =     InInt  "Since which year?";
int cnt = 0;

"\nPapers of ", name, " since ", since, ":\n";
[ SELECT name <= Author && since <= Year;
  cnt = cnt + 1;
  Year, "\t", Title, "\n";
]
"\n", name, " published ", cnt, "papers.\n";
```

U. Kastens: Construction of  
Application Generators  
Using Eli,  
Workshop on Compiler  
Techniques for Application  
Domain Languages ...,  
Linköping, April 1996

# Programming languages as source or target languages

## Programming languages as source languages:

- **Program analysis**  
call graphs, control-flow graph, data dependencies, e. g. for the year 2000 problem
- **Recognition of structures and patterns**  
e. g. for Reengineering

## Program languages as target languages:

- **Specifications (SDL, OMT, UML)**
- **graphic modeling of structures**
- **DSL, Application generator**

**=> Compiler task: Source-to-source compilation**

# Semester project as running example

## A Structure Generator

We are going to develop a tool that implements **record structures**. In particular, the structure generator takes a set of **record descriptions**. Each specifies a **set of named and typed fields**. For each record a **Java class** declaration is to be generated. It contains a constructor method and access methods for the specified record fields.

The tool will be used in an environment where field description are created by other tools, which for example analyze texts for the occurrence of certain phrases. Hence, the descriptions of fields may occur in arbitrary order, and the same field may be described more than once. The structure generator **accumulates the field descriptions** such that for each record a single class declaration is generated which has all fields of that record.

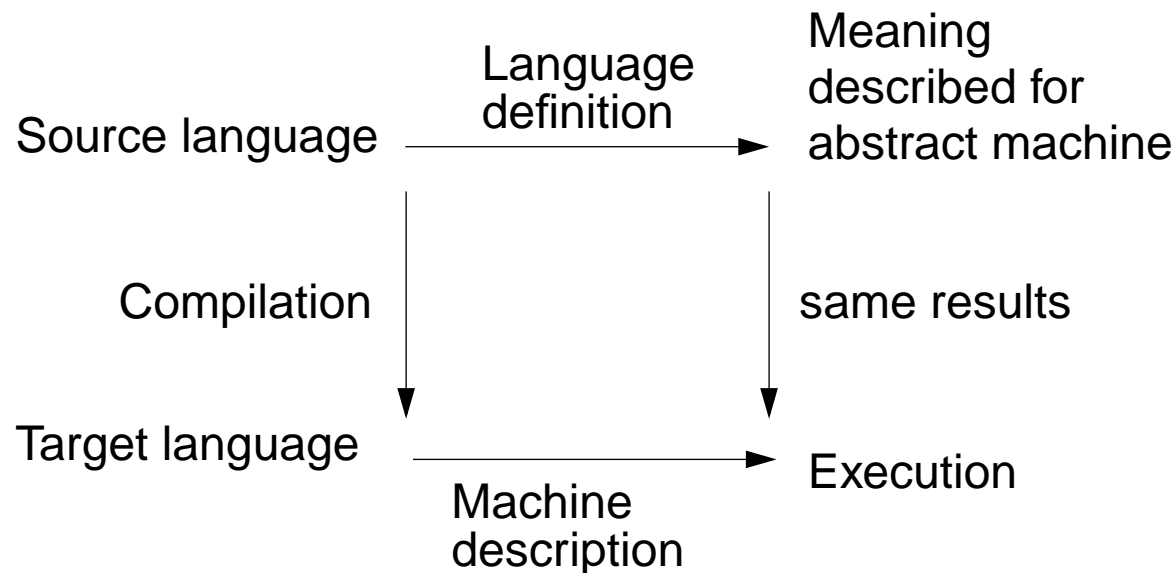
Design a **domain specific language**.

Implement an **application generator** for it.

Apply all **techniques of the course** that are useful for the task.

# Meaning preserving transformation

A **compiler** transforms correct sentences of its **source language** into sentences of its **target language** such that their **meaning is unchanged**.



A **meaning** is defined only for **correct** programs. Compiler task: Error handling

The compiler analyses **static** properties of the program at **compile time**,  
e. g. definitions of Variables, types of expressions. Decides: Is the program **compilable**?

**Dynamic** properties of the program are checked at **runtime**,  
e. g. indexing of arrays. Decides: Is the program **executable**?

But in Java: Compilation of bytecode at runtime, just in time compilation (JIT)



# Example: Tokens and structure

## Character sequence

```
int count = 0; double sum = 0.0; while (count<maxVect) { sum = sum+vect[count]; count++;}
```

## Tokens

```
int count = 0; double sum = 0.0; while (count<maxVect) { sum = sum+vect[count]; count++;}
```

**Expressions**

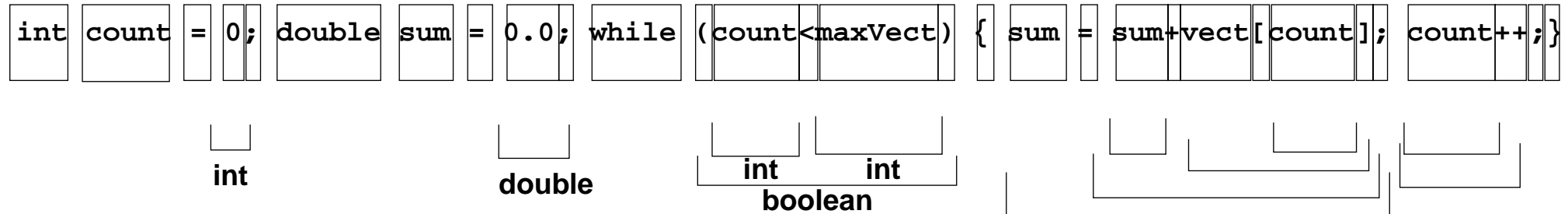
**Declarations**

**Statements**

## Structure

# Example: Names, types, generated code

## Tokens



k1: (count, local variable, int)  
k2: (sum, local variable, double)

k3: (maxVect, member variable, int)  
k4: (vect, member variable, double array)

...

## Names and types

## generated Bytecode

```
0  iconst_0
1  istore_1
2  dconst_0
3  dstore_2
4  goto 19
7  dload_2
8  getstatic #5 <vect[]>
11 iload_1
```

```
12 faload
13 f2d
14 dadd
15 dstore_2
16 iinc 1 1
19 iload_1
20 getstatic #4 <maxVect>
23 if_icmplt 7
```

# Language definition - Compiler task

- **Notation of tokens**

keywords, identifiers, literals

formal definition: regular expressions

**lexical analysis**

- **Syntactic structure**

formal definition: context-free grammar

**syntactic analysis**

- **Static semantics**

binding names to program objects, typing rules  
usually defined by informal texts

**semantic analysis, transformation**

- **Dynamic semantics**

semantics, effect of the execution of constructs  
usually defined by informal texts  
in terms of an abstract machine

**transformation, code generation**

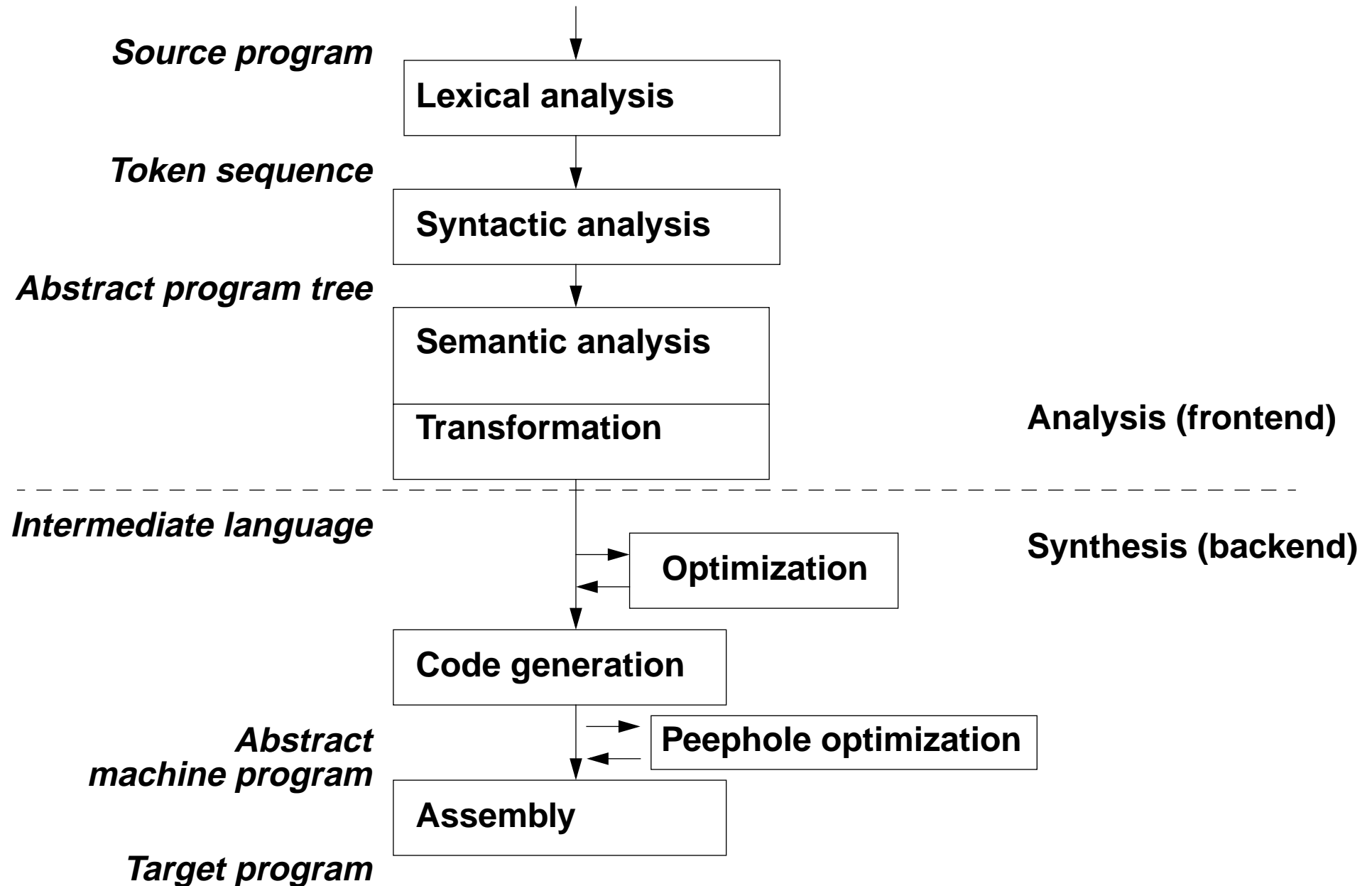
- **Definition of the target language (machine)**

**transformation, code generation  
assembly**

# Compiler tasks

<b>Structuring</b>	<b>Lexical analysis</b>	<b>Scanning</b> <b>Conversion</b>
	<b>Syntactic analysis</b>	<b>Parsing</b> <b>Tree construction</b>
<b>Translation</b>	<b>Semantic analysis</b>	<b>Name analysis</b> <b>Type analysis</b>
	<b>Transformation</b>	<b>Data mapping</b> <b>Action mapping</b>
<b>Encoding</b>	<b>Code generation</b>	<b>Execution-order</b> <b>Register allocation</b> <b>Instruction selection</b>
	<b>Assembly</b>	<b>Instruction encoding</b> <b>Internal Addressing</b> <b>External Addressing</b>

# Compiler structure and interfaces



# Software qualities of the compiler

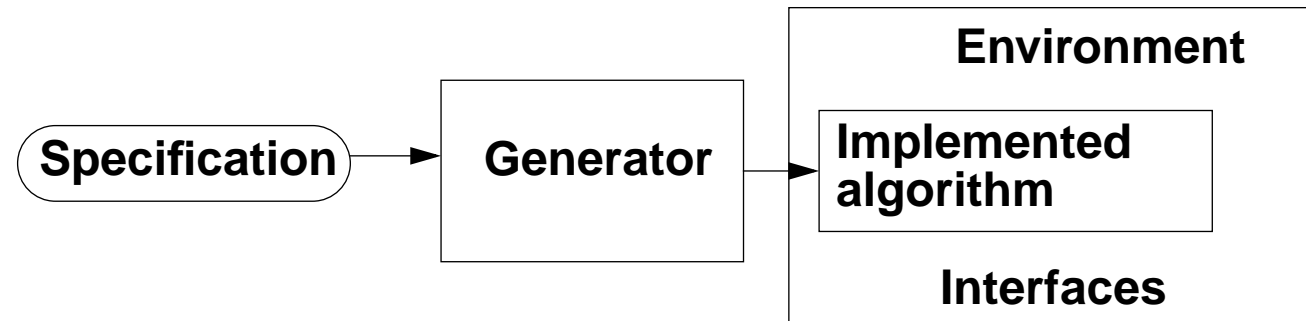
- **Correctness** Translate correct programs correctly.  
Reject wrong programs and give error messages
- **Efficiency** Storage and time used by the compiler
- **Code efficiency** Storage and time used by the generated code  
Compiler task: Optimization
- **User support** Compiler task: Error handling  
(recognition, message, recovery)
- **Robustness** Give a reasonable reaction on every input

# Strategies for compiler construction

- Obey exactly to the language definition
- Use generating tools
- Use standard components
- Apply standard methods
- Validate the compiler against a test suite
- Verify components of the compiler

# Generators

**Pattern:**



**Typical compiler tasks solved by generators:**

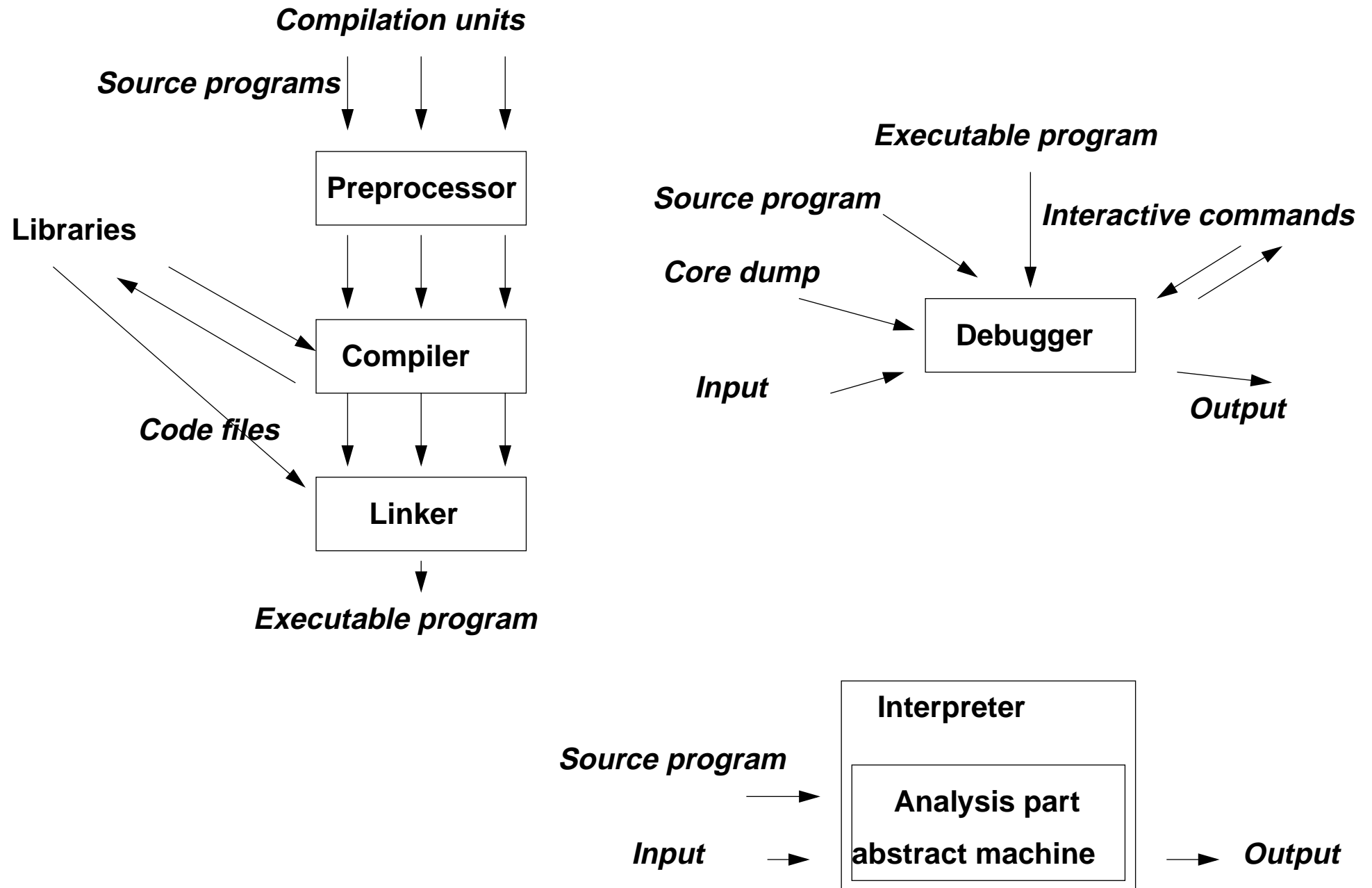
Regular expressions	<b>Scanner generator</b>	Finite automaton
Context-free grammar	<b>Parser generator</b>	Stack automaton
Attribute grammar	<b>Attribute evaluator generator</b>	Tree walking algorithm
Code patterns	<b>Code selection generator</b>	Pattern matching

**integrated system Eli:**





# Environment of compilers



# Compilation and interpretation of Java programs

