

5. Binding Names to Entities

Names in the source code represent **entities** to describe the meaning of the text.

Occurrences of names are **bound to entities**.

Scope rules of the language specify how names are to be bound. E.g.:

- Every name **a**, used as a structure name or as a type name is bound to the same entity.
- A type name **a** is an **applied occurrence** of a name. There must be a **defining occurrences** of **a** somewhere in the text.
- Field names are bound separately for every structure.

some occurrences of names:

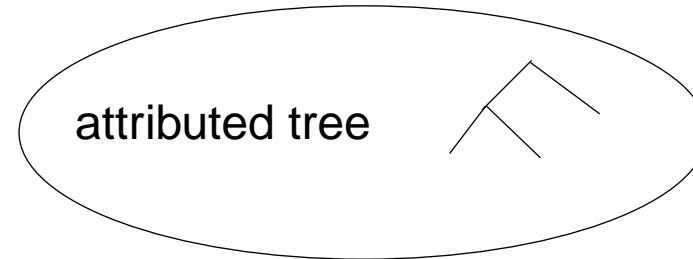
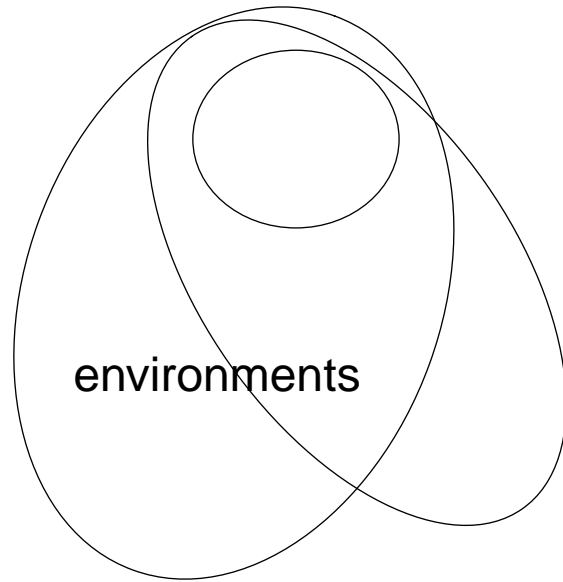
some bindings:

some entities:

```
Customer ( addr: Address;
           account:int;
         )
Address ( name: String;
          zip: int;
         )
Article ( name: String;
          price: int;
         )
```

- a structure (named **Address**)
- a field (named **name**)
- a Structur (named **Article**)
- a different field (named **name**)
- ...

Keys and Properties

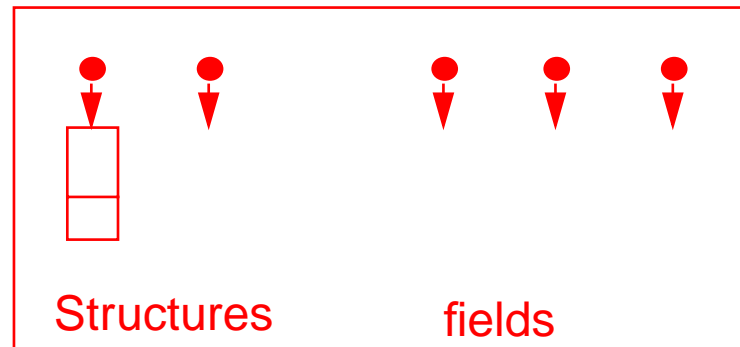


Eli tools implement properties of entities and of environments

**Entities are represented by keys.
Properties are associated to them.**

Structures have a property called **Environment**

Definition module



Entities and their keys

their properties

Bindings and Environments

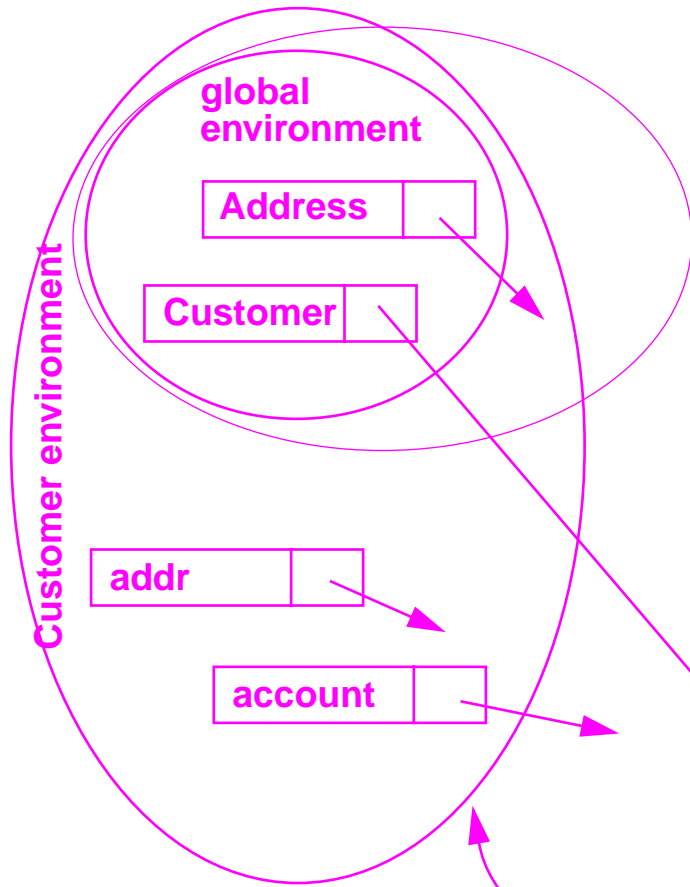
Environment: nested sets of bindings

Binding: associates a name with a key

The **global environment** binds all structure and type names.

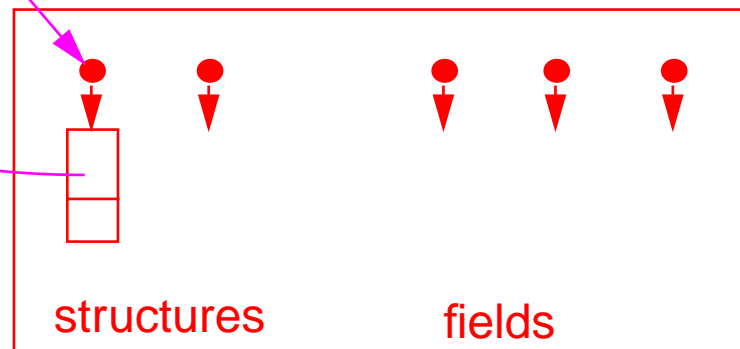
The **environment of a structure** binds its field names.

Eli tools implement properties of entities and of environments



nested environments
sets of bindings

Definition module



Entities and
their keys

their
properties

Attributed Tree for Name Analysis

Attributes of the tree nodes

describe properties of the program construct

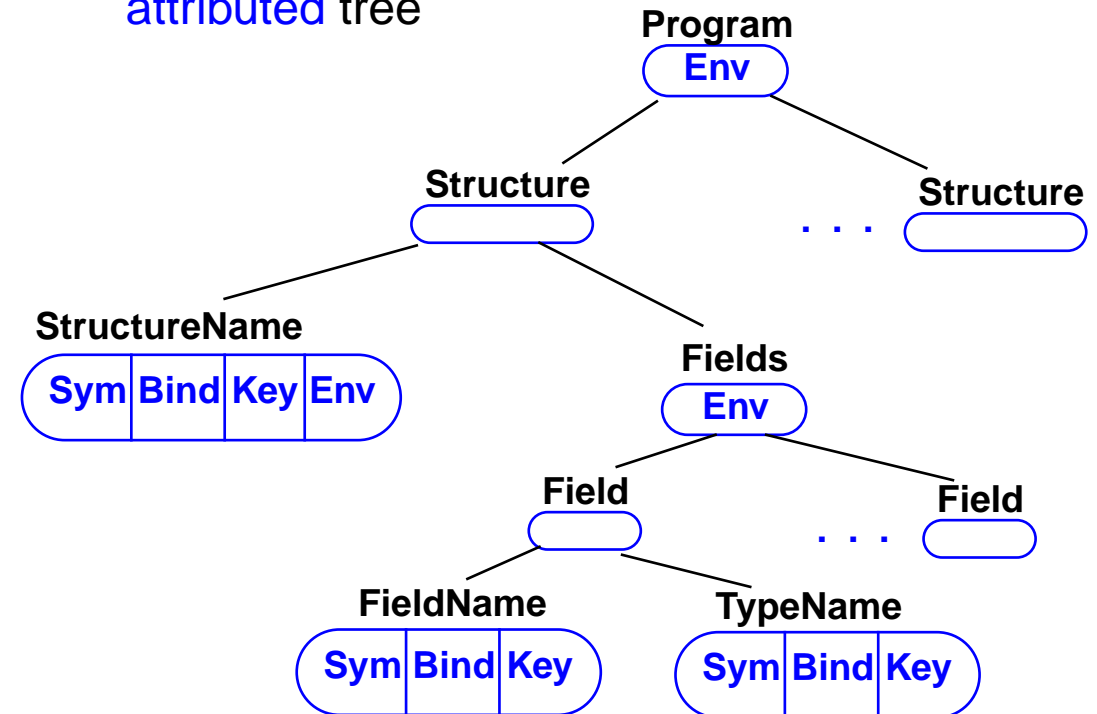
Program has the **global environment**

StructureName and Fields have the **environment of the structure**

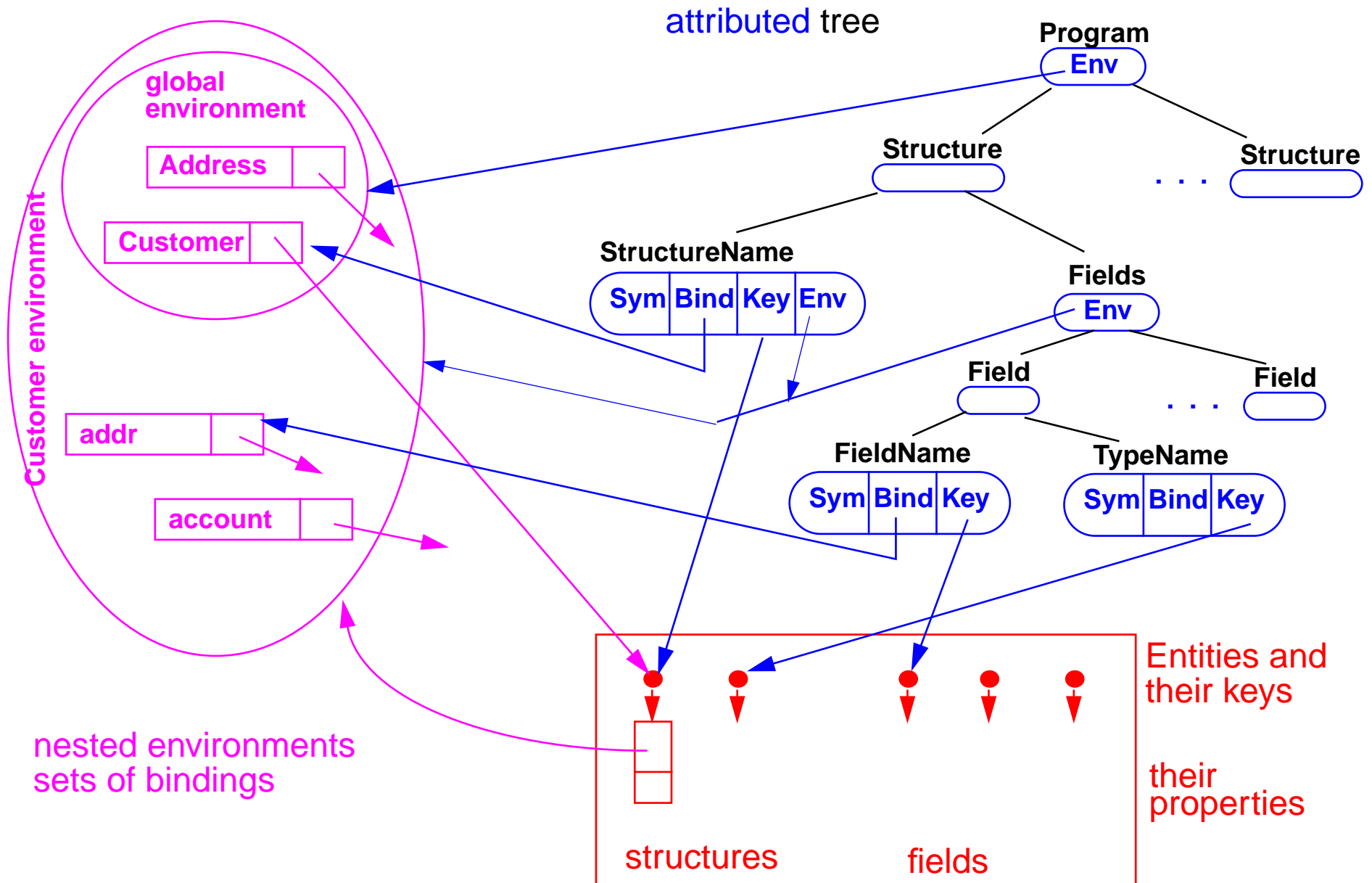
Every node for a name occurrences has attributes for

- the code of the identifier,
- the **binding** of its name, and
- its **key**

attributed tree



Attributes, Environments, and Keys



Environment Module

Implements the abstract data type **Environment**:

hierarchally nested sets (tree) of **bindings (name, environment, key)**

Functions:

NewEnv ()

creates a new environment e , that is the root of a new tree;
used in **root context**

NewScope (e_1)

creates a new environment e_2 that is nested in e_1 .
Every binding of e_1 is a binding of e_2 , too, if it is not hidden
by a binding established for the same name in e_2 ;
used in **range context**

BindIdn (e , id)

creates a new binding (id , e , k), if e does not yet have a
binding for id ; k is then a new key for a new entity;
the result is in both cases the binding (id , e , k);
used for **defining occurrences**.

BindingInEnv (e , id)

yields a binding (id , e_1 , k) of e oder of a surrounding
environment of e ; if there is no such binding it yields NoBinding;
used for **applied occurrences**

BindingInScope (e , id)

yields a binding (id , e , k) of e , if e directly contains such a
binding; NoBinding otherwise; e.g. used for **qualified names**

Example: Names and Entities for the Structure Generator

Abstract syntax

```
RULE: Descriptions LISTOF Import | Structure END;
RULE: Import ::= 'import' ImportNames 'from' FileName END;
RULE: ImportNames LISTOF ImportName END;
RULE: Structure ::= StructureName '(' Fields ')' END;
RULE: Fields LISTOF Field END;
RULE: Field ::= FileName ':' TypeName ';' END;
RULE: StructureName ::= Ident END;
RULE: ImportName ::= Ident END;
RULE: FileName ::= Ident END;
RULE: TypeName ::= Ident END;
```

Different nonterminals for identifiers in different roles,
because different computations are expected, e.g. for
defining and applied occurrences.

Computation of Environment Attributes

Root of the environment hierarchy

```
SYMBOL Descriptions INHERITS RootScope END;
```

Fields play the role of a **Range**.

```
SYMBOL Fields INHERITS RangeScope END;
```

The inherited computation of **Env** is overridden.

```
RULE: Structure ::= StructureName '(' Fields ')'  
COMPUTE  
    Fields.Env = StructureName.Env;  
END;
```

Each structure entity has an **environment as its property**.

```
SYMBOL StructureName COMPUTE  
    SYNT.GotEnvir =  
        IF (EQ (GetEnvir (THIS.Key, NoEnv), NoEnv),  
            ResetEnvir  
                (THIS.Key,  
                    NewScope (INCLUDING Range.Env)));
```

It is **created only once** for every occurrence of a structure entity.

That environment is **embedded in the global environment**.

```
    SYNT.Env =  
        GetEnvir (THIS.Key, NoEnv) <- SYNT.GotEnvir;  
END;
```

In that environment the field names are bound.

Defining and Applied Occurrences of Identifiers

Computations
IdentOcc for all
identifier occurrences.

```
CLASS SYMBOL IdentOcc: Sym: int,  
CLASS SYMBOL IdentOcc COMPUTE  
    SYNT.Sym = TERM;  
END;
```

All **defining** occurrences
bind their names in the
next enclosing Range

```
SYMBOL StructureName  
    INHERITS IdentOcc, IdDefScope END;  
SYMBOL ImportName  
    INHERITS IdentOcc, IdDefScope END;  
SYMBOL FieldName  
    INHERITS IdentOcc, IdDefScope END;
```

Bind an applied
occurrence of an
identifier in the enclosing
environment;
report an error if there is
no valid binding.

```
SYMBOL TypeName  
    INHERITS IdentOcc, IdUseEnv, ChkIdScope END;
```