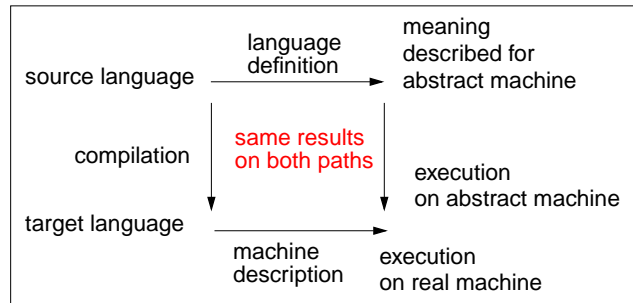


## 1. Language properties - compiler tasks

### Meaning preserving transformation

A **compiler** transforms **any correct sentence** of its **source language** into a sentence of its **target language** such that its **meaning is unchanged**.



A **meaning** is defined only for **all correct** programs => compiler task: error handling

**Static language** properties are analyzed at **compile time**, e. g. definitions of Variables, types of expressions; => determine the transformation, if the program **compilable**

**Dynamic** properties of the program are determined and checked at **runtime**, e. g. indexing of arrays => determine the effect, if the program **executable** (However, just-in-time compilation for Java: bytecode is compiled at runtime.)

## Lecture Programming Languages and Compilers WS 2010/11 / Slide 101

### Objectives:

Understand fundamental notions of compilation

### In the lecture:

The topics on the slide are explained. Examples are given.

- Explain the role of the arcs in the commuting diagram.
- Distinguish compile time and run-time concepts.
- Discuss examples.

## Levels of language properties - compiler tasks

- |  |   |
|--|---|
| <ul style="list-style-type: none"> <li>• <b>a. Notation of tokens</b><br/>keywords, identifiers, literals<br/>formal definition: <b>regular expressions</b></li> </ul>   | <b>lexical analysis</b>                             |
| <ul style="list-style-type: none"> <li>• <b>b. Syntactic structure</b><br/>formal definition: <b>context-free grammar</b></li> </ul>   | <b>syntactic analysis</b>                           |
| <ul style="list-style-type: none"> <li>• <b>c. Static semantics</b><br/>binding names to program objects, typing rules<br/>usually defined by informal texts,<br/>formal definition: <b>attribute grammar</b></li> </ul>   | <b>semantic analysis, transformation</b>            |
| <ul style="list-style-type: none"> <li>• <b>d. Dynamic semantics</b><br/>semantics, effect of the execution of constructs<br/>usually defined by informal texts<br/>in terms of an abstract machine,<br/>formal definition: <b>denotational semantics</b></li> </ul> | <b>transformation, code generation</b>              |
| <b>Definition of target language (target machine)</b>  | <b>transformation, code generation<br/>assembly</b> |

## Lecture Programming Languages and Compilers WS 2010/11 / Slide 102

### Objectives:

Relate language properties to levels of definitions

### In the lecture:

- These are prerequisites of the course "Grundlagen der Programmiersprachen" (see course material GPS-1.16, GPS-1.17).
- Discuss the examples of the following slides under these categories.

### Suggested reading:

Kastens / Übersetzerbau, Section 1.2

### Assignments:

- Exercise 1 Let the compiler produce error messages for each level.
- Exercise 2 Relate concrete language properties to these levels.

### Questions:

Some language properties can be defined on different levels. Discuss the following for hypothetical languages:

- "Parameters may not be of array type." Syntax or static semantics?
- "The index range of an array may not be empty." Static or dynamic semantics?

## Example: Tokens and structure

### Character sequence

```
int count = 0; double sum = 0.0; while (count < maxVect) { sum = sum + vect[count]; count++; }
```

### Tokens

```
int count = 0; double sum = 0.0; while (count < maxVect) { sum = sum + vect[count]; count++; }
```

Expressions

Declarations

Statements

### Structure

## Lecture Programming Languages and Compilers WS 2010/11 / Slide 103

### Objectives:

Get an idea of the structuring task

### In the lecture:

Some requirements for recognizing tokens and deriving the program structure are discussed along the example:

- kinds of tokens,
- characters between tokens,
- nested structure

### Questions:

Where do you find the exact requirements for the structuring tasks?

## Example: Names, types, generated code

```
int count = 0; double sum = 0.0; while (count < maxVect) { sum = sum + vect[count]; count++; }
```

### Structure

int

double

int int

boolean

k1: (count, local variable, int)

k2: (sum, local variable, double)

k3: (maxVect, member variable, int) ...

k4: (vect, member variable, double array)

Static properties: names and types

### generated Bytecode

```
0 iconst_0          12 faload
1 istore_1          13 f2d
2 dconst_0          14 dadd
3 dstore_2          15 dstore_2
4 goto 19           16 iinc 1 1
7 dload_2           19 iload_1
8 getstatic #5 <vect[]> 20 getstatic #4 <maxVect>
11 iload_1           23 if_icmplt 7
```

## Lecture Programming Languages and Compilers WS 2010/11 / Slide 104

### Objectives:

Get an idea of the name analysis and transformation task

### In the lecture:

Some requirements for these tasks are discussed along the example:

- program objects and their properties,
- program constructs and their types
- target program

### Questions:

- Why is the name (e.g. count) a property of a program object (e.g. k1)?
- Can you impose some structure on the target code?

## Compiler tasks

Structuring	Lexical analysis	Scanning Conversion
	Syntactic analysis	Parsing Tree construction
Translation	Semantic analysis	Name analysis Type analysis
	Transformation	Data mapping Action mapping
Encoding	Code generation	Execution-order Register allocation Instruction selection
	Assembly	Instruction encoding Internal Addressing External Addressing

## Lecture Programming Languages and Compilers WS 2010/11 / Slide 105

### Objectives:

Language properties lead to decomposed compiler tasks

### In the lecture:

- Explain tasks of the rightmost column.
- Relate the tasks to chapters of the course.

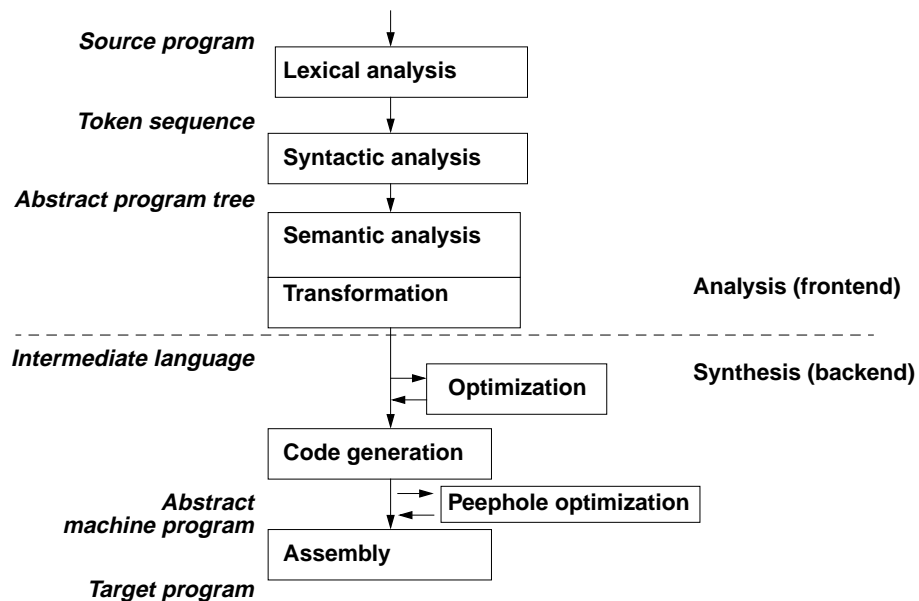
### Suggested reading:

Kastens / Übersetzerbau, Section 2.1

### Assignments:

Learn the German translations of the technical terms.

## Compiler structure and interfaces



## Lecture Programming Languages and Compilers WS 2010/11 / Slide 106

### Objectives:

Derive compiler modules from tasks

### In the lecture:

In this course we focus on the analysis phase (frontend).

### Suggested reading:

Kastens / Übersetzerbau, Section 2.1

### Assignments:

Compare this slide with [U-08](#) and learn the translations of the technical terms used here.

### Questions:

Use this information to explain the example on slides PLaC-1.3, 1.4

## Software qualities of the compiler

PLaC-1.7

- **Correctness**      Compiler translates correct programs correctly; rejects wrong programs and gives error messages
- **Efficiency**        Storage and time used by the compiler
- **Code efficiency**    Storage and time used by the generated code; compiler task: optimization
- **User support**      Compiler task: Error handling (recognition, message, recovery)
- **Robustness**        Compiler gives a reasonable reaction on every input; does not break on any program

## Lecture Programming Languages and Compilers WS 2010/11 / Slide 107

### Objectives:

Consider compiler as a software product

### In the lecture:

Give examples for the qualities.

### Questions:

Explain: For a compiler the requirements are specified much more precisely than for other software products.

## Strategies for compiler construction

PLaC-1.8

- **Obey exactly to the language definition**
- **Use generating tools**
- **Use standard components**
- **Apply standard methods**
- **Validate the compiler against a test suite**
- **Verify components of the compiler**

## Lecture Programming Languages and Compilers WS 2010/11 / Slide 108

### Objectives:

Apply software methods for compiler construction

### In the lecture:

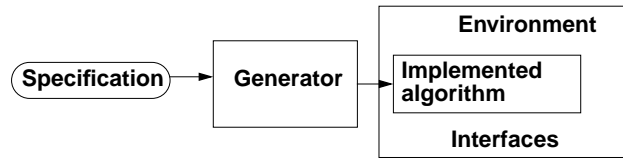
It is explained that effective construction methods exist especially for compilers.

### Questions:

What do the specifications of the compiler tasks contribute to more systematic compiler construction?

## Generate from specifications

Pattern:



### Typical compiler tasks solved by generators:

Regular expressions	<b>Scanner generator</b>	Finite automaton
Context-free grammar	<b>Parser generator</b>	Stack automaton
Attribute grammar	<b>Attribute evaluator generator</b>	Tree walking algorithm
Code patterns	<b>Code selection generator</b>	Pattern matching

integrated system Eli:



## Lecture Programming Languages and Compilers WS 2010/11 / Slide 109

### Objectives:

Usage of generators in compiler construction

### In the lecture:

The topics on the slide are explained. Examples are given.

### Suggested reading:

Kastens / Übersetzerbau, Section 2.5

### Assignments:

- [Exercise 5](#): Find as many generators as possible in the Eli system.

## Compiler Frameworks (Selection)

### Amsterdam Compiler Kit: (Uni Amsterdam)

The Amsterdam Compiler Kit is fast, lightweight and retargetable compiler suite and toolchain written by Andrew Tanenbaum and Criel Jacobs. Intermediate language EM, set of frontends and backends

### ANTLR: (Terence Parr, Uni San Francisco)

ANother Tool for Language Recognition, (formerly PCCTS) is a language tool that provides a framework for constructing recognizers, compilers, and translators from grammatical descriptions containing Java, C#, C++, or Python actions

### CoCo: (Uni Linz)

Coco/R is a compiler generator, which takes an attributed grammar of a source language and generates a scanner and a parser for this language. The scanner works as a deterministic finite automaton. The parser uses recursive descent.

### Eli: (Unis Boulder, Paderborn, Sydney)

Combines a variety of standard tools that implement powerful compiler construction strategies into a domain-specific programming environment called Eli. Using this environment, one can automatically generate complete language implementations from application-oriented specifications.

### SUIF: (Uni Stanford)

The SUIF 2 compiler infrastructure project is co-funded by DARPA and NSF. It is a free infrastructure designed to support collaborative research in optimizing and parallelizing compilers.

## Lecture Programming Languages and Compilers WS 2010/11 / Slide 109a

### Objectives:

General information on compiler tool kits

### In the lecture:

Some characteristics of the systems are explained.

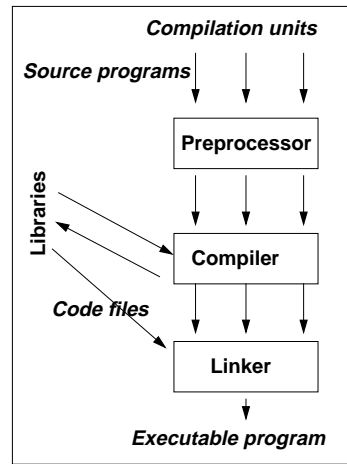
### Suggested reading:

Kastens / Übersetzerbau, Section 2.5

### Assignments:

- Find more information on the system in the Web

## Environment of compilers



**Preprocessor** cpp substitutes text macros in source programs, e.g.

```
#include <stdio.h>
#include "module.h"

#define SIZE 32
#define SEL(ptr, fld) ((ptr)->fld)
```

Separate compilation of compilation units

- with interface specification, consistency checks, and language specific linker: Modula, Ada, Java
- without ...; checks deferred to system linker: C, C++

## Lecture Programming Languages and Compilers WS 2010/11 / Slide 110

### Objectives:

Understand the cooperation between compilers and other language tools

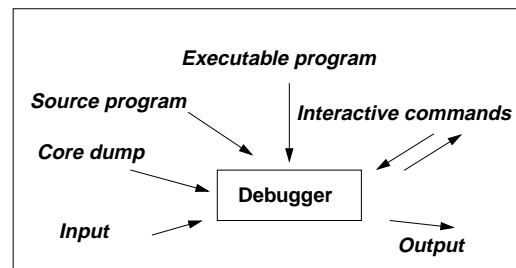
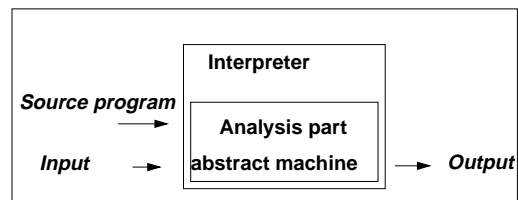
### In the lecture:

- Explain the roles of language tools
- Explain the flow of information

### Suggested reading:

Kastens / Übersetzerbau, Section 2.4

## Interpreter and Debugger



## Lecture Programming Languages and Compilers WS 2010/11 / Slide 110a

### Objectives:

Understand the cooperation between compilers and other language tools

### In the lecture:

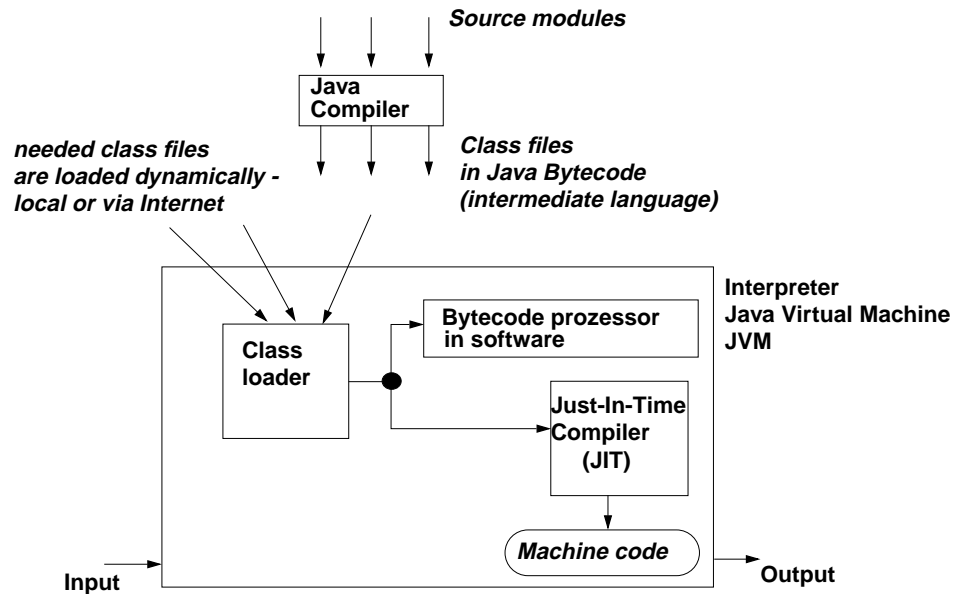
- Explain the roles of language tools
- Explain the flow of information

### Suggested reading:

Kastens / Übersetzerbau, Section 2.4

# Compilation and interpretation of Java programs

PLaC-1.11



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## Lecture Programming Languages and Compilers WS 2010/11 / Slide 111

### Objectives:

Special situation for Java

### In the lecture:

Explain the role of the abstract machine JVM:

- Interpretation of bytecode.
- JIT: Compiles and optimizes while executing the program.
- JVM: Loads class files while executing the program.

### Questions:

- explain why the JVM can not rely on the type checks made by a compiler.