

Construction of Attribute Evaluators

For a given attribute grammar an attribute evaluator is constructed:

- It is **applicable to any tree** that obeys the abstract syntax specified in the rules of the AG.
- It performs a **tree walk** and **executes computations** in visited contexts.
- The execution order obeys the **attribute dependences**.

Pass-oriented strategies for the tree walk: **AG class:**

k times **depth-first left-to-right**

k times **depth-first right-to-left**

alternatingly left-to-right / right-to left

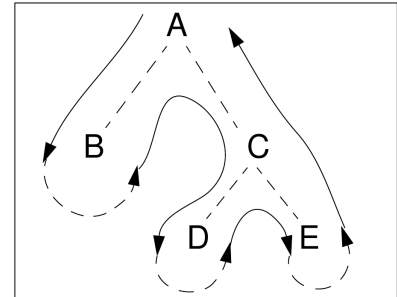
once **bottom-up (synth. attributes only)**

LAG (k)

RAG (k)

AAG (k)

SAG



AG is checked if attribute dependences

fit to desired pass-oriented strategy; see LAG(k) check.

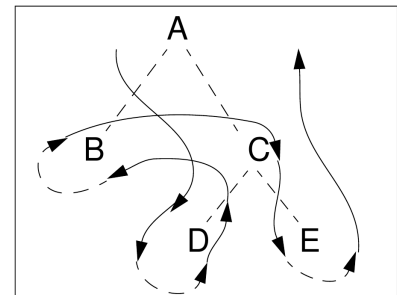
non-pass-oriented strategies:

visit-sequences:

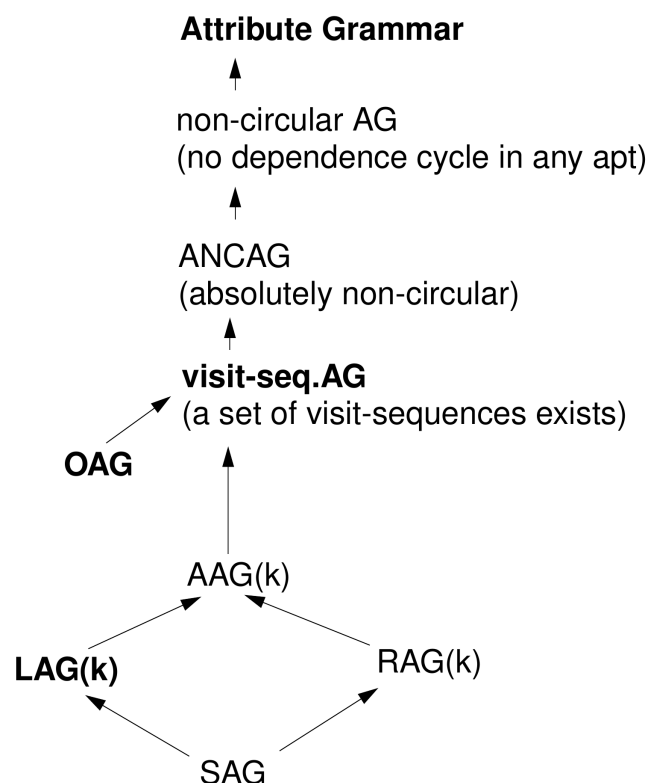
an individual plan for each rule of the abstract syntax

OAG

A generator fits the plans to the dependences of the AG.



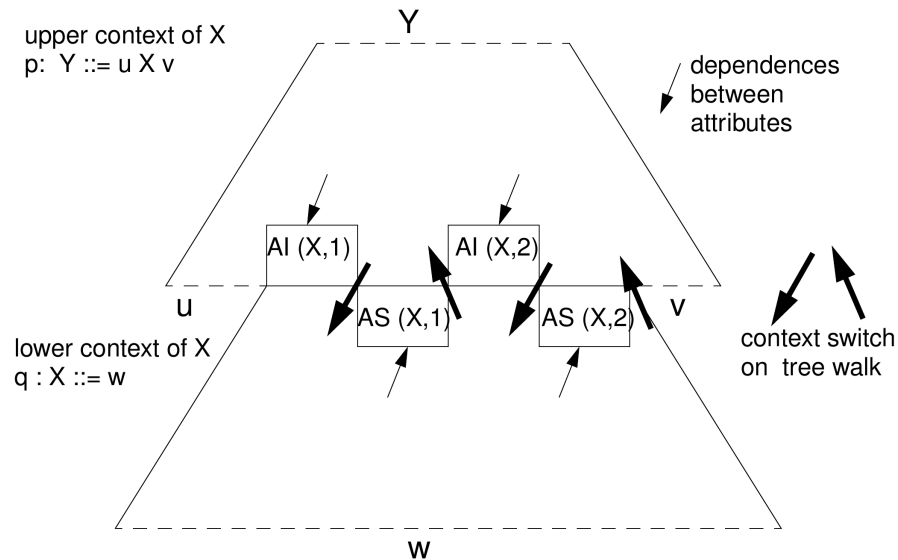
Hierarchy of AG Classes



The sets $AI(X)$ and $AS(X)$ are **partitioned** each such that

$AI(X, i)$ is computed before the i -th visit of X

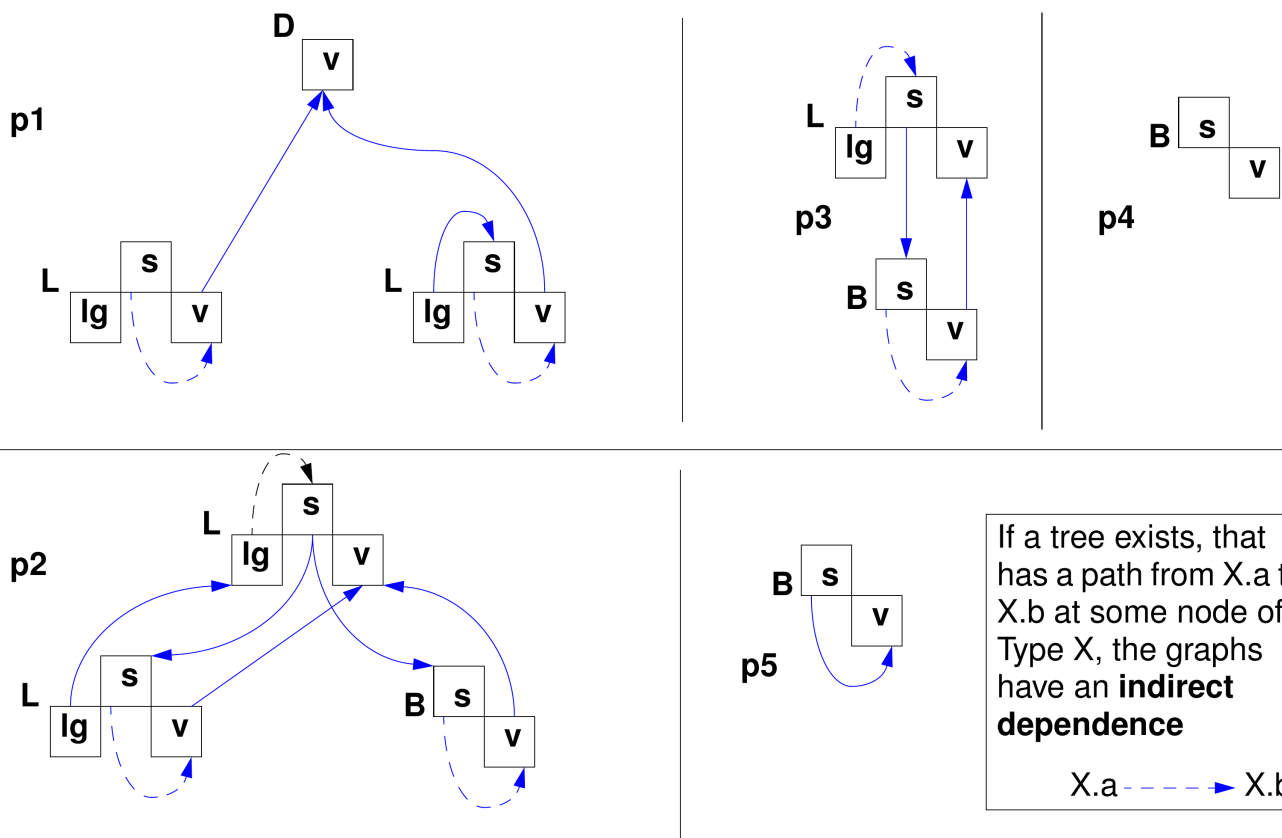
$AS(X, i)$ is computed during the i -th visit of X



Necessary precondition for the existence of such a partition:

No node in any tree has direct or indirect dependencies that contradict the evaluation order of the sequence of sets: $AI(X, 1)$, $AS(X, 1)$, ..., $AI(X, k)$, $AS(X, k)$

Dependency Analysis for AG "Binary Numbers"



Pass-Oriented Evaluation: LAG(k)

An AG is a LAG(k), if:

For each symbol X there is an **attribute partition** $A(X, 1), \dots, A(X, k)$, such that the attributes in $A(X, i)$ can be computed in the i -th depth-first left-to-right pass.

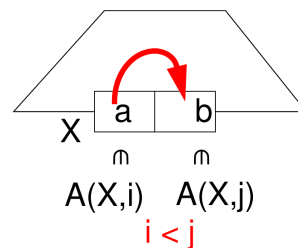
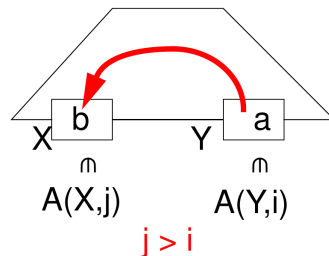
Crucial dependences:

In every dependence graph every dependence

- $Y.a \rightarrow X.b$ where X and Y occur on the **right-hand side** and Y is **right of** X implies that **$Y.a$ belongs to an earlier pass than $X.b$** , and
- $X.a \rightarrow X.b$ where X occurs on the **right-hand side** implies that **$X.a$ belongs to an earlier pass than $X.b$**

Necessary and sufficient condition over dependence graphs - expressed graphically:

A dependency
from right to left



A dependence
at one symbol
on the right-hand
side

LAG (k) Algorithm

Algorithm checks whether there is a $k \geq 1$ such that an AG is LAG(k).

Method:

compute iteratively $A(1), \dots, A(k)$;
in each iteration try to allocate all remaining attributes to the current pass, i.e. $A(i)$;
remove those which can not be evaluated in that pass

Algorithm:

Set $i=1$ and $\text{Cand} =$ all attributes

repeat

set $A(i) = \text{Cand}$; set Cand to empty;

while still attributes can be removed from $A(i)$ do

remove an attribute $x.b$ from $A(i)$ and add it to Cand if

- there is a **crucial dependence**

$Y.a \rightarrow X.b$ s.t.

X and Y are on the right-hand side, Y to the right of X and $Y.a$ in $A(i)$ or

$X.a \rightarrow X.b$ s.t. X is on the right-hand side and $X.a$ is in $A(i)$

- $x.b$ depends on an attribute that is not yet in any $A(i)$

if Cand is empty:

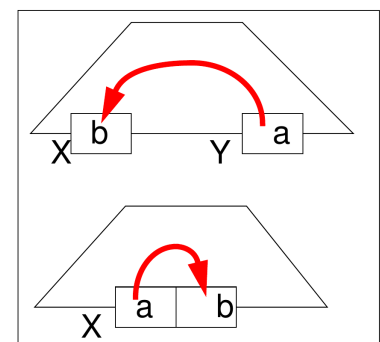
exit: the AG is **LAG(k)** and all attributes are assigned to their passes

if $A(i)$ is empty:

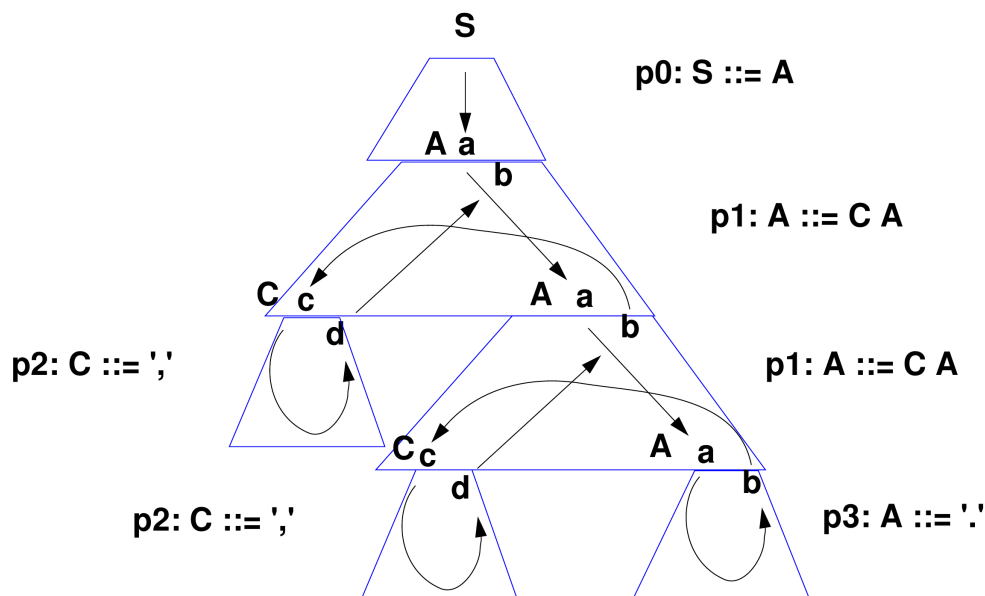
exit: the AG is **not LAG(k) for any k**

else:

set $i = i + 1$



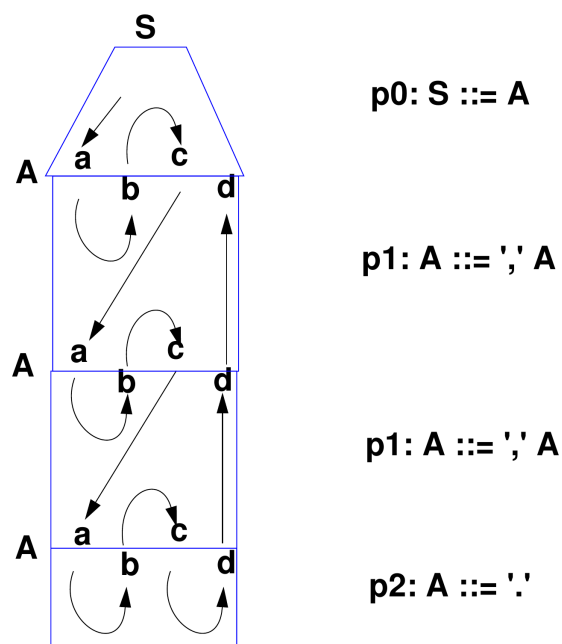
AG not LAG(k) for any k



**A.a can be allocated to the first left-to-right pass.
C.c, C.d, A.b can not be allocated to any pass.**

The AG is RAG(1), AAG(2) and can be evaluated by visit-sequences.

AG not evaluable in Passes



No attribute can be allocated to any pass for any strategy.

The AG can be evaluated by visit-sequences.