## Parallel Programming WS 2014/2015 - Solution 5

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## **Solution for Exercise 4**

Original loop:

```
for I = 1 to 3 do
  for J = 1 to 3 do
        A[I, J] = A[I - 1, J + 1] + 1
        endfor
endfor
```

a) The iteration space including a dependence vector.

```
x x x x
x x x
x x x
D = 1
-1
```

b) It is illegal to permute the I and J loops because the resulting dependence vector is negative:

c) Applying a skewing transformation with factor 1:

```
T = 1 0
1 1
T * D = 1 0 * 1 = 1
1 1 - 1 0
```

Legal.

d) The resulting iteration space including a dependence vector:

```
x coord (3,6)
x x
x x x
x x x
x coord (1,2)
D = 1
0
```

e) Permuting the I and J loops of the resulting code:

Legal.

f) The resulting iteration space including a dependence vector:

```
 \begin{array}{c} x & x & x \\ x & x & x \\ x & x & x \\ \end{array}  coord (4,3) - coord(6,3)
 \begin{array}{c} x & x & x \\ x & x & x \\ \end{array}  coord (2,1) - coord(4,1)
 D = 0 \\ 1 \end{array}
```

g) The loop bounds of the transformed loop derived from the drawing:

i' runs from 2 to 6
j' runs from max(1, i' - 3) to min(3, i' - 1)

h) Homework: Deriving the loop bounds mathematically:

The inverse transformation matrix T<sup>-1</sup>:

0 1 1 -1

The bounds equation derived from the original program is

```
B * i <= -1

j 3

-1

3

==>

-1 0 * i <= -1

1 0 j 3

0 -1 -1

0 1 3
```

The bounds equation derived for the transformed program is

 $B * T^{-1} * i' <= -1$   $j' \qquad 3$  -1 3==> 0 -1 \* i' <= -1  $0 1 j' \qquad 3$   $-1 1 \qquad -1$   $1 -1 \qquad 3$ which yields j' >= 1 j' <= 3 j' - i' <= -1 i' - j' <= 3

```
==>
```

2 <= i' <= 6 j' >= max (1, i' - 3) j' y= min (3, i' - 1) which corresponds to the bounds we derived graphically. Additionally we compute the transformed loop body:

The old iteration variables in terms of the new ones:

T<sup>-1</sup> \* i' = i j' j ==> 0 1 \* i' = i 1 -1 j' j ==> i = j' j = i' - j'

Together the transformed program looks as follows:

```
for i' = 2 to 6 do
  for j' = max (1, i' - 3) to min (3, i' - 1) do
        A[j', i' - j'] = A[j' - 1, i' - j' + 1] + 1
    endfor
endfor
```

It computes the array elements in the order

A[11], A[12], A[21], A[13], A[22], A[31], A[23], A[32], A[33]

## **Solution for Exercise 5**

```
for I = 0 to N do
  for J = 0 to M do
        A[I + 1, J] = A[I, J - 1] + A[I + 1, J - 2]
      endfor
endfor
```

These are the necessary steps:

• 1. Draw the iteration space.

A rectangle from (0,0) to (3,3):

- x x x x x x x x x x x x x x x x
- 2. Compute the dependence vectors and draw examples of them into the iteration space.

The dependence vectors are  $(1, 1)^{T}$  and  $(0, 2)^{T}$ .

- 3. Apply a skewing transformation with factor 1 and draw the iteration space.
- 4. Apply a permutation transformation and draw the iteration space.
- 5. Compute the matrix of the composed transformation and use it to transform the dependence vectors.

```
Skewing matrix S: 1 0

1 1

Permutation matrix: 0 1

1 0

T = P * S: 1 1

1 0

Transformed dependence vectors T * D:

1 1 1 0 2 2

* =

1 0 1 2 1 0
```

• 6. Compute the inverse of the transformation matrix and use it to transform the index expressions.

Inverse Transformation Matrix T<sup>-1</sup>:

0 1 1 -1

Transformation of index expressions:

0 1 ip i \* = 1 -1 jp j

yields i = jp and j = ip - jp.

• 7. Specify the loop bounds by inequalities and transform them by the inverse of the transformation matrix.

 $B * T^{-1} * IP \le c$ 

-1	0						0
1	0		0 1		ip		N
		*		*		<=	
0	-1		1 -1		jp		0
0	1						М

simplifies to

0 ·	-1				0
0	1		ip		Ν
		*		<=	
-1	1		jp		0
1	-1				М

such that

0 <= ip <= M + N max(0, ip - M) <= jp <= min (ip, N)

• 8. Write the complete loops with new loop variables ip and jp and new loop bounds.

```
for I = 0 to M+N do
  for J = max(0, I-M) to min(I, N) do
        A[J + 1, I - J] = A[J, I - J - 1] + A[J + 1, I - J - 2]
   endfor
endfor
```

If N = M = 2 the transformed loop executes the computations:

These are the same computations that are executed by the original loop. The computations of the inner loop are parallelizable.